Problem 1: Textbook RSA and hybrid encryption

A common variant of textbook RSA is the following: During key generation, the modulus $N$ is chosen as usual. We chose $e$ as $e := 3$ (instead of random). Then $d$ is chosen with $ed \equiv 1 \mod \varphi(N)$ (as usual). This is implemented by the Python functions `rsa_keygen`, `rsa_enc`, `rsa_dec` below.

We use this in a “hybrid encryption”, which first picks an AES key $k$, encrypts it with RSA, and then encrypts the actual message with AES using the key $k$. (Functions `hyb_enc`, `hyb_dec`.)

Your task is to write an adversary that, given the public key $pk$, and the hybrid encryption $c$ of some message $m$, finds $m$. That is, fill in the function body of the function `adv` below so that the function `test_adv` prints `Success. The adversary broke the scheme.

**Hint:** We discuss/discussed in the practice the problem with RSA with $e = 3$ when RSA-encrypting short messages.

(You find the following file on the lecture webpage, too.)

```python
#!/usr/bin/python3
# Use "pip install sympy" (possibly with sudo) to install sympy
# And "Crypto" might need "pip install pycrypto" if it’s not installed
import sympy, math, Crypto, random

prime_len = 1024

def egcd(a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, y, x = egcd(b % a, a)
        return (g, x - (b // a) * y, y)
def modinv(a, m):
    g, x, y = egcd(a, m)
    if g != 1:
        raise Exception('modular inverse does not exist')
```
else:
    return x % m

def rsa_keygen():
    while True:
        try:
            p = sympy.ntheory.generate.randprime(2**prime_len,2**(prime_len+1))
            q = sympy.ntheory.generate.randprime(2**prime_len,2**(prime_len+1))
            e = 3
            N = p*q
            phiN = (p-1)*(q-1)
            pk=(N,e)
            sk=(N,modinv(e,phiN))
            return (pk,sk)
        except Exception as e:
            pass

# Rough ad-hoc algorithm, not optimized

def exp_mod(a,e,N):
    res = 1
    b = a
    i = 0
    while e>=2**i: # Invariant: b=a**(2**i)
        if e & 2**i != 0:
            e -= 2**i
            res = (res*b) % N
            b=(b*b) % N
            i += 1
    assert e==0
    return res

# Just a test
assert exp_mod(23123,323,657238293) == ((23123**323) % 657238293)

def rsa_enc(pk,m):
    (N,e) = pk
    return exp_mod(m,e,N)

def rsa_dec(sk,c):
    (N,d) = sk
    return exp_mod(c,d,N)

def int_to_bytes(i,len): # Not optimized
res = []
for j in range(len):
    res.append(i%256)
    i = i>>8
return bytes(res)

def aes_cbc_enc(k,m):
    from Crypto.Cipher import AES
    from Crypto import Random
    assert len(m)%AES.block_size == 0
    k = int_to_bytes(k,AES.block_size)
    iv = Random.new().read(AES.block_size)
    cipher = AES.new(k, AES.MODE_CBC, iv)
    return iv + cipher.encrypt(m)

def aes_cbc_dec(k,m):
    from Crypto.Cipher import AES
    from Crypto import Random
    k = int_to_bytes(k,AES.block_size)
    iv = m[:AES.block_size]
    cipher = AES.new(k, AES.MODE_CBC, iv)
    return cipher.decrypt(m[AES.block_size:]),

# Just a test
assert aes_cbc_dec(2123414234,aes_cbc_enc(2123414234,b'hello there test')) == b'hello there test'

def hyb_enc(pk,m):
    assert isinstance(m,bytes)
    k = random.getrandbits(256)
    aes_k_m = aes_cbc_enc(k,m)
    assert m == aes_cbc_dec(k,aes_k_m)
    rsa_pk_k = rsa_enc(pk,k)
    return (rsa_pk_k,aes_k_m)

def hyb_dec(sk,c):
    (c1,c2) = c
    k = rsa_dec(sk,c1)
    m = aes_cbc_dec(k,c2)
    return m

def adv(pk,c):
    m = b"put the right message here"
    return m

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Problem 2: Malleability of textbook RSA

The adversary get a textbook RSA encryption \( c = E(pk, m) \) for some unknown message \( m \). The adversary also knows \( pk = (N, e) \). The adversary wants to compute \( c' = E(pk, 2m) \). (This is a specific example of malleability.) How can the adversary efficiently compute \( c' \) from \( c \) and \( pk \)?

You may assume that \( 0 \leq m < N/2 \).

Problem 3: Security proofs (bonus problem)

Recall the definition of IND-OT-CPA [Definition 3 in the lecture notes]. There, we defined security by saying that if the adversary cannot distinguish between an encryption of \( m_0 \) and \( m_1 \) in the sense that it will output 1 will almost the same probability in both cases.

Consider the following variant of the definition:

**Definition 1 (IND-OT-CPA – variant)** An encryption scheme \((KG, E, D)\) is \((\tau, \varepsilon)\)-
**IND-OT-CPA’ secure** if for any \(\tau\)-time algorithm \(A\) we have that

\[
\left| \Pr[b^* = b : b \leftarrow \{0, 1\}, k \leftarrow KG(), (m_0, m_1) \leftarrow A(), c \leftarrow E(k, m_b), b^* \leftarrow A(c)] - \frac{1}{2} \right| \leq \varepsilon.
\]

(Here we quantify only over algorithms \(A\) that output \((m_0, m_1)\) with \(|m_0| = |m_1|\).)

That is, the message \(m_b\) is encrypted (with \(b\) random), and we want that the adversary cannot guess \(b\) with probability much different from \(\frac{1}{2}\). (Guessing with probability \(\frac{1}{2}\) is always possible since \(b\) is just a single bit.)

We wish to prove that if \((KG, E, D)\) is \((\tau, \varepsilon)\)-IND-OT-CPA’ secure, then \((KG, E, D)\) is \((\tau, 2\varepsilon)\)-IND-OT-CPA secure.

**Note:** The converse also holds, but we will not prove that.

(a) Assume an adversary \(A\) that breaks \((\tau, 2\varepsilon)\)-IND-OT-CPA security. Let

\[
\alpha_0 := \Pr[b^* = 1 : k \leftarrow KG(), (m_0, m_1) \leftarrow A(), c \leftarrow E(k, m_0), b^* \leftarrow A(c)]
\]

and

\[
\alpha_1 := \Pr[b^* = 1 : k \leftarrow KG(), (m_0, m_1) \leftarrow A(), c \leftarrow E(k, m_1), b^* \leftarrow A(c)].
\]

What do we know about \(\alpha_0\) and \(\alpha_1\) (by definition of IND-OT-CPA and the fact that \(A\) breaks IND-OT-CPA)?

(b) Compute

\[
\beta := \left| \Pr[b^* = b : k \leftarrow KG(), b \leftarrow \{0, 1\}, (m_0, m_1) \leftarrow A(), c \leftarrow E(k, m_b), b^* \leftarrow A(c)] - \frac{1}{2} \right|.
\]

(As a formula using \(\alpha_0\) and \(\alpha_1\).)

(c) Using (a) and (b), show that if \(A\) breaks \((\tau, 2\varepsilon)\)-IND-OT-CPA, then \(A\) breaks \((\tau, \varepsilon)\)-IND-OT-CPA’. (Hence: IND-OT-CPA’ implies IND-OT-CPA.)

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\footnote{This is not an established name!}