Functional Programming IO Monad

Jevgeni Kabanov

Department of Computer Science University of Tartu

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Outline









Haskell is a *pure* and *lazy* language

Pure

- Purity means that every function result is uniquely determined by its parameters
- Compiler is free to inline any function as it pleases
- Runtime can cache any function call

Lazy

- Laziness means that values are computed on need
- If they are not needed, they are not computed
- Runtime is free to precompute them or reorder computations

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The following function should read a character from standard input

getChar::Char

Let's see an example:

get2chars = [getChar, getChar]

- How many characters will be read?
- In what order will they be read?

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Problem 1

Let's try to solve these problems:

 $getChar::Int
ightarrow Char \\ get2chars = [getChar 1, getChar 2]$

Now the values are distinct, yet the order is still undefined

Problem 2

Also we now need to add the same *Int* argument to *get2chars* for the same reason:

```
get4chars::Int 
ightarrow [Char]get4chars=get2chars1++get2chars2
```

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```
get4chars::Int 
ightarrow [Char]get4chars = get2chars 1 + get2chars 2
```

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```

Problem 3

getChar::Int
ightarrow Charget2chars = [getChar 1, getChar 2]

• What will be called first?

- What we really need is some kind of dependency!
- The only kind of runtime dependency Haskell provides is value dependency:

 $egin{aligned} getchar :: Int
ightarrow (Char, Int) \ get2chars _ = [\,a, b\,] extbf{where} (\,a, i) = getChar \, 1 \ (b, _) = getChar \, i \end{aligned}$

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Problem 4

The problem now is get2chars – the parameter is neither used nor unique and compiler might still reorder or omit it.

get4chars = [get2chars 1, get2chars 2]

What we need is a unique parameter associated with *every* IO action:

get2chars :: Int
ightarrow (String, Int) get2chars i0 = ([a, b], i2) where (a, i1) = getChar i0 (b, i2) = getChar i1 get4chars i0 = (a + b, i2) where (a, i1) = get2chars i0(b, i2) = get2chars i1

Problem 4

The problem now is get2chars – the parameter is neither used nor unique and compiler might still reorder or omit it.

get4chars = [get2chars 1, get2chars 2]

What we need is a unique parameter associated with *every* IO action:

$$get2chars :: Int \rightarrow (String, Int)$$

 $get2chars i0 = ([a, b], i2)$ where $(a, i1) = getChar i0$
 $(b, i2) = getChar i1$
 $get4chars i0 = (a + b, i2)$ where $(a, i1) = get2chars i0$
 $(b, i2) = get2chars i1$

Problem 5

However there is still one problem left, can you see it?:

$$get2chars i0 = ([a, b], i2)$$
 where $(a, i1) = getChar i0$
 $(b, i2) = getChar i0$

A simple mistake and our carefully built up IO is ruined!

NB!

What we need is a way for compiler to ensure that the value is unique!

Problem 5

However there is still one problem left, can you see it?:

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When Haskell was created three main ways for doing pure, lazy IO were known.

- Input-output streams
- Ontinuations
- Threaded unique value

Of course Haskellers chose the fourth way!

Clean

Interestingly enough a Haskell-like language Clean chose to introduce *uniqueness typing* that made the third way possible directly

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Real World IO

We continue the previous example by interpreting the threaded value as the token of the world state:

```
getChar :: RealWorld \rightarrow (Char, RealWorld)
```

In that case the program is of type

 $main :: RealWorld \rightarrow ((), RealWorld)$

We can introduce a type synonym and rewrite it as:

 $f{type} \ IO \ a = RealWorld
ightarrow (a, RealWorld) \ main :: IO () \ getChar :: IO \ Char$

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type IO $a = RealWorld \rightarrow (a, RealWorld)$ main :: IO () getChar :: IO Char

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Hello, Monad!

Let's see a program reading two characters from input:

$$\begin{array}{l} main \ world0 = \mathbf{let} \ (a, world1) = getChar \ world0\\ (b, world2) = getChar \ world1\\ \mathbf{in} \ ([a, b], world2) \end{array}$$

We could hide the token from the programmer by threading it automatically:

 $(\gg=) :: IO \ a \rightarrow (a \rightarrow IO \ b) \rightarrow IO \ b$ $(action1 \gg= action2) \ world0 =$ $let (a, world1) = action1 \ world0$ $(b, world2) = action2 \ a \ world1$ in (b, world2)

main =

 $getChar \gg = \lambda a
ightarrow getChar \gg = \lambda b
ightarrow ([a, b], w))$

Hello, Monad!

Let's see a program reading two characters from input:

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(b, world2) = getChar world1
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 $getChar \gg = \lambda a \rightarrow getChar \gg = \lambda b \rightarrow (\lambda w \rightarrow ([a, b], w))$

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Doing It with Class

Now we can make IO a a monad!

instance Monad (IO a) where

$$(\gg=) :: IO \ a \rightarrow (a \rightarrow IO \ b) \rightarrow IO \ b$$

 $(action1 \gg= action2) \ world0 =$
let $(a, world1) = action1 \ world0$
 $(b, world2) = action2 \ a \ world1$
in $(b, world2)$
return :: $a \rightarrow IO \ a$
return $a = (\lambda world \rightarrow (a, world))$

We can now use the *do*-notation so the previous example becomes

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 $egin{array}{l} main = {f do} \ a \leftarrow getChar \ b \leftarrow getChar \ return \ [a, b] \end{array}$

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$$egin{array}{l} main = {f do} \ a \leftarrow getChar \ b \leftarrow getChar \ return \ [a, b] \end{array}$$

Reflections

- It is logical to make *IO* a abstract, so that programmer would be protected from making mistakes.
- Then a value of *IO* a is nothing more than an *action* or a *computation* to be executed.
- However this also means that all IO functions must be predefined with priority access to IO monad.

IO functions

IO module necessarily contains functions for working with:

- Files and file systems
- Mutable variables
- Random numbers
- System clock

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Console Operations

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Definition

Most of console functions are self-explanatory:

```
putChar :: Char \rightarrow IO ()
putStr :: String \rightarrow IO ()
putStrLn :: String \rightarrow IO ()
print :: Show a \Rightarrow a \rightarrow IO ()
readLn :: Read a \Rightarrow IO a
getChar :: IO Char
getLine :: IO String
getContents :: IO String
```

Console Operations

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Example

IO Actions as Values

Example

Let's define a list of IO actions:

```
ioActions :: [IO ()]
ioActions =
 [(print "Hello!"),
 (putStr "just kidding"),
 (getChar >> return ())]
```

Remember, that IO $a = RealWorld \rightarrow (a, RealWorld)$, so these are usual functional values.

IO Actions as Values

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Example

However these IO actions can also be used directly:

```
main = do
head ioActions
ioActions !! 1
last ioActions
```

Output:

```
"Hello!"
just kidding(Wait for input)
```

IO Action Lists

Definition

Sequencing a list of actions can be reused (foldr is just a reminder):

$$\begin{array}{ll} foldr & :: (a \rightarrow b \rightarrow b) \rightarrow b \rightarrow [a] \rightarrow b \\ foldr \ f \ z \ [] = & z \\ foldr \ f \ z \ (x:xs) = f \ x \ (foldr \ f \ z \ xs) \\ sequence _ :: Monad \ m \Rightarrow [m \ a] \rightarrow m \ () \\ sequence _ = foldr \ (\gg) \ (return \ ()) \\ \end{array}$$

A lot more interesting monad combinators are available and we will examine them later.

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IO Action Lists

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Definition

The previous example then becomes:

```
ioActions :: [IO ()]
ioActions =
  [(print "Hello!"),
  (putStr "just kidding"),
  (getChar >> return ())]
main = do sequence_ ioActions
```

Output:

"Hello!"
just kidding(Wait for input)

Definition

Files are opened as *Handles* and can be used accordingly:

```
openFile :: String \rightarrow IOMode \rightarrow IO Handle

hSeek :: Handle \rightarrow SeekMode \rightarrow Integer \rightarrow IO ()

hGetChar :: Handle \rightarrow IO Char

hPutChar :: Handle \rightarrow Char \rightarrow IO ()

hClose :: Handle \rightarrow IO ()

data SeekMode

= AbsoluteSeek

| RelativeSeek

| SeekFromEnd
```

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Example

Let's define a function reading one character from a given handle and position:

```
readi \ h \ i = do
hSeek \ h \ i \ AbsoluteSeek
hGetChar \ h
```

Next we want to open the file and just read characters from it:

```
egin{aligned} 	ext{readfilei} :: String 	o IO \ (Integer 	o IO \ Char) \ 	ext{readfilei} name = 	extbf{do} \ h \leftarrow openFile \ name \ ReadMode \ return \ (readi \ h) \end{aligned}
```

Example

Let's define a function reading one character from a given handle and position:

```
readi h i = do
hSeek h i AbsoluteSeek
hGetChar h
```

Next we want to open the file and just read characters from it:

```
readfilei :: String 
ightarrow IO (Integer 
ightarrow IO Char) \ readfilei name = do \ h \leftarrow openFile name ReadMode \ return (readi h)
```

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Example

Since IO functions can be higher order we can apply *readi* repeatedly:

IORef

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Definition

IORef a allows to create mutable variables:

```
data IORef a

newIORef :: a \rightarrow IO (IORef a)

readIORef :: IORef a \rightarrow IO a

writeIORef :: IORef a \rightarrow a \rightarrow IO ()

modifyIORef :: IORef a \rightarrow (a \rightarrow a) \rightarrow IO ()
```

IORef

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Example

```
egin{aligned} main &= \mathbf{do} \ varA \leftarrow newIORef \ 0 \ a0 \leftarrow readIORef \ varA \ writeIORef \ varA \ 1 \ a1 \leftarrow readIORef \ varA \ print \ (a0, a1) \end{aligned}
```

Haskell Objects

Figure

Let's try to define a Figure object:

We will define only one figure – a circle:

 $circle :: Point \rightarrow Radius \rightarrow IO \ Figure$ $type \ Point = (Int, Int) \ -- \ point \ coordinates$ $type \ Radius = Int \ -- \ circle \ radius \ in \ points$

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Haskell Objects

Drawing

We can test the objects when we define them with this code:

```
main = do
  figures \leftarrow sequence [circle (10, 10) 5,
    rectangle (10, 10) (20, 20)
  drawAll figures
  mapM (\lambda fig \rightarrow move fig (10, 10)) figures
  drawAll figures
drawAll :: [Figure] \rightarrow IO ()
drawAll \ figures = do
  putStrLn "Drawing figures:"
  mapM draw figures
```

Circle

```
circle center radius = do
  centerVar \leftarrow newIORefcenter
 let drawF = do
    center \leftarrow readIORef centerVar
    putStrLn (" Circle at " ++ show center
       ++ " with radius " ++ show radius)
 let moveF (addX, addY) = do
    (x, y) \leftarrow readIORef centerVar
    writeIORef centerVar (x + addX, y + addY)
  return \ Figure { draw = drawF, move = moveF }
```

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Rectangle

```
rectangle from to = do
  fromVar \leftarrow newIORef from
  to Var \leftarrow newIORef to
  let drawF = do
    from \leftarrow readIORef from Var
    to \leftarrow readIORef \ to Var
    putStrLn (" Rectangle " ++ show from
       ++ "-" ++ show to
  let moveF(addX, addY) = do
    (from X, from Y) \leftarrow readIORef from Var
    (toX, toY) \leftarrow readIORef \ toVar
    writeIORef from Var (from X + addX, from Y + addY)
    writeIORef to Var (toX + addX, toY + addY)
  return  Figure { draw = drawF, move = moveF }
```

Higher Order IO

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Lazy IO

How lazy is the IO monad?

```
list1, list2 :: IO [Int]
list1 = return \circ repeat \$ 0
list2 = sequence \circ repeat \circ return \$ 0
main = do
lst \leftarrow list1
putStrLn \circ show \circ take 10 \$ lst
readLn
lst2 \leftarrow list2
putStrLn \circ show \circ take 10 \$ lst2
```

Higher Order IO

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```
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putStrLn \circ show \circ take 10 \$ lst
readLn
lst2 \leftarrow list2
putStrLn \circ show \circ take 10 \$ lst2
```

The first list will show the first 10 elements, while the second one will go infinite.

Advantages

- IO in Haskell is pure and lazy
- The order of operations that end up executed is guaranteed by the compiler
- IO actions are treated in Haskell first-class, they can be passed around, combined and curried

- Every new function that needs to do side-effects needs access to IO monad implementation
- IO actions tend to contaminate pure code, lifting it to IO monad because some particular parts need to do IO
- Since IO operations are the fastest, a large portion of a program may end up under IO monad

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