

Functional Programming

QuickCheck: Automated Random Testing

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Random Testing

Outline

- Testing is very important in programming
- In JUnit and alike we collect test cases that are prearranged argument-result pairs
- In Haskell there is HUnit which does the same
- However we could do better
 - Since functions are pure we can test them against properties
 - Since data types are structural we can try generating random data samples
- Random testing enjoys some of the benefits of formal verification without nearly as much pain!

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reverse examples

Property

We concatenated in a wrong order:

```
propRevApp2 :: [Int] → [Int] → Bool  
propRevApp2 xs ys =  
  reverse (xs ++ ys) ≡ reverse ys ++ reverse xs
```

Output

```
Test> quickCheck propRevApp2  
OK, passed 100 tests.
```

reverse examples

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Output

```
Test> quickCheck propRevApp2  
OK, passed 100 tests.
```

reverse examples

Property

Let's check if you can reverse before concatenating:

```
propRevApp1 :: [Int] → [Int] → Bool  
propRevApp1 xs ys =  
    reverse (xs ++ ys) ≡ reverse xs ++ reverse ys
```

Output

```
Test> quickCheck propRevApp1
```

Falsifiable, after 4 tests:

```
[-3,-4,-4]
```

```
[-4,-1,1,1]
```

reverse examples

Property

Let's check if you can reverse before concatenating:

```
propRevApp1 :: [Int] → [Int] → Bool  
propRevApp1 xs ys =  
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Falsifiable, after 4 tests:

```
[-3,-4,-4]
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Distribution examples

Property

The following property asserts that addition and multiplication distribute:

```
propDistributiveI :: Int → Int → Int → Bool  
propDistributiveI a b c =  
    a * (b + c) ≡ (a * b) + (a * c)
```

Output

```
Test> propDistributiveI  
OK, passed 100 tests.
```

Distribution examples

Property

The following property asserts that addition and multiplication distribute:

```
propDistributiveI :: Int → Int → Int → Bool  
propDistributiveI a b c =  
    a * (b + c) ≡ (a * b) + (a * c)
```

Output

```
Test> propDistributiveI  
OK, passed 100 tests.
```

Distribution examples

Property

The same property for *Floats* fails:

```
propDistributiveF :: Float → Float → Float → Bool
propDistributiveF a b c =
  a * (b + c) ≡ (a * b) + (a * c)
```

Output

```
Test> quickCheck propDistributiveF
```

Falsifiable, after 7 tests:

3.0

-2.666667

3.75

Distribution examples

Property

The same property for *Floats* fails:

```
propDistributiveF :: Float → Float → Float → Bool
propDistributiveF a b c =
  a * (b + c) ≡ (a * b) + (a * c)
```

Output

```
Test> quickCheck propDistributiveF
```

Falsifiable, after 7 tests:

3.0

-2.666667

3.75

insert and ordered

Definition

For the next several slides we will consider a function which inserts an element into an ordered list.

```
insert e (x : xs) =  
  if e < x then e : x : xs else x : (insert e xs)  
insert e [] = [e]
```

ordered tests whether the list is ordered:

```
ordered :: Ord a ⇒ [a] → Bool  
ordered [] = True  
ordered (x : []) = True  
ordered (x1 : x2 : xs) =  
  if x1 ≤ x2 then ordered (x2 : xs)  
  else False
```

insert examples

Property

We would want to test whether *insert* works, but this has point only on ordered lists:

```
propInsert1 :: Int → [Int] → Bool
propInsert1 x xs =
  if ordered xs
    then ordered (insert x xs)
    else True
```

Since QuickCheck does not work on polymorphic types we choose *Ints* here.

Output

```
Test> propInsert1
OK, passed 100 tests.
```

insert examples

Property

We would want to test whether *insert* works, but this has point only on ordered lists:

```
propInsert1 :: Int → [Int] → Bool
propInsert1 x xs =
  if ordered xs
    then ordered (insert x xs)
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```

Since QuickCheck does not work on polymorphic types we choose *Ints* here.

Output

```
Test> propInsert1
OK, passed 100 tests.
```

insert examples

Property

But did this actually tell us anything? How do we know how many lists were ordered?

```
propInsert2 :: Int → [Int] → Property
propInsert2 x xs =
  (length xs ≡ 0 ∨ ¬ (ordered xs)) 'trivial'
  if ordered xs
  then ordered (insert x xs)
  else True
```

Output

```
*Test> quickCheck propInsert2
OK, passed 100 tests (82% trivial).
```

insert examples

Property

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```
propInsert2 :: Int → [Int] → Property
propInsert2 x xs =
  (length xs ≡ 0 ∨ ¬ (ordered xs)) 'trivial'
  if ordered xs
  then ordered (insert x xs)
  else True
```

Output

```
*Test> quickCheck propInsert2
OK, passed 100 tests (82% trivial).
```

insert examples

Property

`==>` is the QuickCheck combinator that makes it test only the fitting values:

```
propInsert3 :: Int → [Int] → Property  
propInsert3 x xs =  
    ordered xs ==> ordered (insert x xs)
```

Output

```
Test> propInsert3  
OK, passed 100 tests.
```

insert examples

Property

`==>` is the QuickCheck combinator that makes it test only the fitting values:

```
propInsert3 :: Int → [Int] → Property  
propInsert3 x xs =  
    ordered xs ==> ordered (insert x xs)
```

Output

```
Test> propInsert3  
OK, passed 100 tests.
```

insert examples

Property

How well do we actually test? Can this pass?

```
insBad a [] = [a]
insBad a y
  | (length y) > 4 = y ++ [a]
  | otherwise = insert a y

propInsertBad1 :: Int → [Int] → Property
propInsertBad1 x xs =
  ordered xs ==> ordered (insBad x xs)
```

Output

```
Test> quickCheck propInsertBad1
OK, passed 100 tests.
```


insert examples

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```
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Output

```
Test> quickCheck propInsertBad1
OK, passed 100 tests.
```

insert examples

Property

```
propInsertBad2 :: Int → [Int] → Property  
propInsertBad2 x xs =  
    ordered xs ==>  
        collect (length xs) $ ordered (insBad x xs)
```

Output

```
Test> quickCheck propInsertBad2
```

```
OK, passed 100 tests.
```

```
53% 0.
```

```
24% 1.
```

```
14% 2.
```

```
8% 3.
```

```
1% 4.
```

insert examples

Property

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propInsertBad2 :: Int → [Int] → Property  
propInsertBad2 x xs =  
  ordered xs ==>  
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53% 0.
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```
14% 2.
```

```
8% 3.
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```
1% 4.
```

insert examples

Property

```
propInsertBad3 :: Int → [Int] → Property
propInsertBad3 x xs =
  ordered xs ==>
    classify (ordered (x : xs)) "at-head" $
    classify (ordered (xs ++ [x])) "at-tail" $
    ordered (insBad x xs)
```

Output

```
Test> quickCheck propInsertBad3
OK, passed 100 tests.
53% at-head, at-tail.
20% at-tail.
20% at-head.
```

insert examples

Property

```
propInsertBad3 :: Int → [Int] → Property
propInsertBad3 x xs =
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    classify (ordered (x : xs)) "at-head" $
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Output

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Test> quickCheck propInsertBad3
OK, passed 100 tests.
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Generators

Outline

- We test mostly trivial or very simple cases (only one insert in the middle of the list!)
- Just checking whether list is ordered is not enough!
- We need a way to generate ordered lists!

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Gen a

Definition

Generators are instances of the Monad class with the (simplified) concrete representation:

newtype *Gen a* = *Gen* (*Rand* \rightarrow *a*)

The types of `bind` and `return` suggest we can use them as combinators to build complex generators out of simpler ones:

return :: *a* \rightarrow *Gen a*

($\gg=$) :: *Gen a* \rightarrow (*a* \rightarrow *Gen b*) \rightarrow *Gen b*

Arbitrary a

Definition

The type class *Arbitrary a* denotes types for which we can generate random values:

```
class Arbitrary a where  
  arbitrary :: Gen a
```

And these values are used in a property by applying *forAll*:

```
forAll :: (Show a, Testable b) =>  
  Gen a -> (a -> b) -> Property
```

Arbitrary instances

Definition

Given a function $choose :: (Int, Int) \rightarrow Gen\ Int$, we write:

```
instance Arbitrary Int where  
  arbitrary = choose (-42, 42)
```

We can use the built-in *liftM2* monad function to add pairs to the *Arbitrary* class.

```
instance (Arbitrary a, Arbitrary b) =>  
  Arbitrary (a, b) where  
    arbitrary = liftM2 (,) arbitrary arbitrary
```

Enumeration generator

Definition

The $\text{oneof} :: [Gen\ a] \rightarrow Gen\ a$ combinator randomly selects one generator from a list. Elements are weighted equally.

```
data Prof = Steve | Stephanie | Benjamin
instance Arbitrary Prof where
    arbitrary = oneof
    [return Steve, return Stephanie, return Benjamin]
```

We can also define $\text{Arbitrary}\ [a]$ using oneof :

```
instance Arbitrary a  $\Rightarrow$  Arbitrary [a] where
    arbitrary = oneof
    [return [], liftM2 (:) arbitrary arbitrary]
```

List generator

Definition

Our previous instantiation of *Arbitrary* $[a]$ created empty lists half the time. To fix this we use

$frequency :: [(Int, Gen a)] \rightarrow Gen a$:

```
instance Arbitrary a  $\Rightarrow$  Arbitrary [a] where
  arbitrary = frequency
    [(1, return []),
     (4, liftM2 (:) arbitrary arbitrary)]
```

Trees generator

Definition

We can also instantiate a tree generator:

```
data Tree a = Leaf a | Branch (Tree a) (Tree a)
instance Arbitrary a =>
  Arbitrary Tree a where
  arbitrary = frequency
    [(1, LiftM Leaf arbitrary),
     (2, LiftM2 Branch arbitrary arbitrary)]
```

What's wrong with this definition?

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```

What's wrong with this definition?

Sized generators

Definition

We can ensure generated data structures have finite size by adding an explicit size parameter to $Gen\ a$. Our definition becomes

$$\text{newtype } Gen\ a = Gen\ (Int \rightarrow Rand \rightarrow a)$$

and is used with a new combinator:

$$sized :: (Int \rightarrow Gen\ a) \rightarrow Gen\ a$$

Tree generator

Definition

The following tree definition will produce a trees with no more elements than the parameter to *arbTree*. Note that this parameter is passed in by *sized* and is a global constant.

```
data Tree a = Leaf a | Branch (Tree a) (Tree a)
instance Arbitrary a =>
  Arbitrary Tree a where
    arbitrary = sized arbTree
    arbTree 0 = liftM Leaf arbitrary
    arbTree n = frequency
      [(1, liftM Leaf arbitrary),
       (2, liftM2 Branch
          (arbTree (n `div` 2))
          (arbTree (n `div` 2)))]
```

insert examples

Definition

Back to our problem:

```
insBad a [] = [a]
insBad a y
  | (length y) > 4 = y ++ [a]
  | otherwise = insert a y

propInsertBad1 :: Int → [Int] → Property
propInsertBad1 x xs =
  ordered xs ==> ordered (insBad x xs)
```

Output

```
Test> quickCheck propInsertBad1
OK, passed 100 tests.
```

orderedList

Definition

Now we can define *orderedList* generator:

```
orderedList = do  
  a ← frequency [(1, return []),  
    (7, liftM2 (:) arbitrary arbitrary)]  
  return (sort a)
```

Example

Definition

And finally fail the example!

```
propInsertBad4 :: Int → Property  
propInsertBad4 x =  
  forAll orderedList $ λxs → ordered (insBad x xs)
```

Output

```
*Test> quickCheck propInsertBad4  
Falsifiable, after 10 tests:  
-6  
[-8,-4,-3,0,5]
```

Infinite Structures

Definition

Infinite structures will cause infinite loops:

```
propDoubleCycle1 :: [Int] → Property  
propDoubleCycle1 xs =  
  ¬ (null xs) ==>  
    cycle xs ≡ cycle (xs ++ xs)
```

Infinite Structures

Definition

However we can control them up to any finite size:

```
propDoubleCycle2 :: [Int] → Int → Property
propDoubleCycle2 xs n =
  ¬ (null xs) ∧ n ≥ 0 ==>
    take n (cycle xs) ≡ take n (cycle (xs ++ xs))
```

Functions

Definition

Let's try to define random functions by throwing away the input and generating a random result. In this case:

$$\begin{aligned} \text{propFunc1} &:: (Int \rightarrow Int) \rightarrow Int \rightarrow Bool \\ \text{propFunc1 } f \ x &= (f \circ (+2)) \ x \equiv (f \circ (*2)) \ x \end{aligned}$$

Output

Test> quickCheck propFunc1 OK, passed 100 tests.

Functions

Outline

- We need a functional dependency between input and output, or we can get wrong results
- Type of $Gen\ (a \rightarrow b)$ is $Int \rightarrow Rand \rightarrow a \rightarrow b$
 - This is equivalent to $a \rightarrow Int \rightarrow Rand \rightarrow b$
 - And $a \rightarrow Gen\ b$
- It's not clear we can make a value of one type into a generator for another.
 - However maybe we can use arbitrary Ints to transform generators with $variant :: Int \rightarrow Gen\ a \rightarrow Gen\ a$.
 - We can certainly make specific types into Ints:

```
coarbitrary b = if b
  then variant 1
  else variant 0
```


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```
coarbitrary b = if b
               then variant 1
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```

Functions

Outline

- In Haskell, the right way to generalize this is with a type class.

```
class Coarbitrary a where  
    coarbitrary :: a → Gen b → Gen b
```

- We then define *Arbitrary* in terms of *Coarbitrary* (and a helper function to match the types).

```
instance (Coarbitrary a, Arbitrary b) ⇒  
    Arbitrary (a → b) where  
    arbitrary =  
        promote (λa → coarbitrary a arbitrary)
```

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```


Functions

Definition

variant :: *Int* → *Gen a* → *Gen a*

variant *v* (*Gen m*) =

Gen ($\lambda n\ r \rightarrow m\ n\ (rands\ r\ !!\ (v + 1))$)

where

rands *r0* = *r1* : *rands* *r2* **where** (*r1*, *r2*) = *split* *r0*

promote :: (*a* → *Gen b*) → *Gen* (*a* → *b*)

promote *f* =

Gen ($\lambda n\ r \rightarrow \lambda a \rightarrow \text{let } Gen\ m = f\ a\ \text{in } m\ n\ r$)

Functions

Definition

```
instance Coarbitrary Bool where
  coarbitrary b =
    if b then variant 0 else variant 1

instance Coarbitrary Int where
  coarbitrary n =
    variant (if n ≥ 0 then 2 * n else 2 * (-n) + 1)

instance Coarbitrary Char where
  coarbitrary c = variant (ord c)
```

Functions

Definition

And back to the example:

$$\begin{aligned} \text{propFunc1} &:: (\text{Int} \rightarrow \text{Int}) \rightarrow \text{Int} \rightarrow \text{Bool} \\ \text{propFunc1 } f \ x &= (f \circ (+2)) \ x \equiv (f \circ (*2)) \ x \end{aligned}$$

Output

```
*Test> quickCheck propFunc1
Falsifiable, after 0 tests:
*function*
-3
```

Implementation

Definition

```
newtype Property = Prop (Gen Result)  
class Testable a where  
    property :: a → Property  
instance Testable Bool where  
    property b = Prop (return $ resultBool b)  
instance Testable Property where  
    property prop = prop  
instance (Arbitrary a, Show a, Testable b) ⇒  
    Testable (a → b) where  
    property f = forAll arbitrary f
```

Testing Monads

Outline

- It is impossible to random test *IO* monad
- *ST* monad can be tested by randomly generating lists of actions
- It is not too comfortable
- However since functions like $==>$ are defined on *Propertys*, we need to redefine them on a monad transformer *PropertyM*
- QuickCheck2 provides support for that

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Shrinking

Outline

- Often we find a counter example, but it's way too big to understand the underlying cause
- In such a case it is possible to start shrinking the example to find a subexample that still causes the function to fail
- This is implemented as an extra function *shrink* in *Arbitrary* class that generates all substructures
- QuickCheck2 implements these and some extra for most common structures

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- This is implemented as an extra function *shrink* in *Arbitrary* class that generates all substructures
- QuickCheck2 implements these and some extra for most common structures