Modelling Cyclic Structures by Nested Datatypes

Varmo Vene

Teooriapäevad, Koke, 3-5 February 2006

Motivation

- Algebraic datatypes provide a nice way to represent tree-like structures.
- Lazy languages, eg. Haskell, allow to build also cyclic structures.

```
cycle = 1 : 2 : cycle
```

 Allows to represent complete infinite structures in finite memory

Motivation

- However, there is no support for manipulating cyclic structures
- Eg. mapping over cyclic list gives an infinite list
 map (+1) cycle ==> [2,3,2,3,2,3,2,3,....
- In fact, there is no way to distinguish cyclic structures from infinite ones
- Our aim is to represent cyclic structures inductively, hence to separate them from infinite (coinductive) structures.
- This gives the ability to explicitly manipulate cyclic structures either directly or using generic operations like fold, etc.

Cyclic Lists – 1st attempt

- Sheard, Fegaras 1996 (??)
- Definition:

Examples:

```
clist1 = Rec (\lambda xs -> Cons 1 xs)
clist2 = Rec (\lambda xs -> Cons 1 (Cons 2 xs))
clist3 = Cons 1 (Rec (\lambda xs -> Cons 2 xs))
```

- Doesn't have unique representation
- Requires higher-order recursive datatypes



Cyclic Lists – 2nd attempt

Definition:

• Examples:

```
clist1 = Cons 1 (Var 1)
clist2 = Cons 1 (Cons 2 (Var 1))
clist3 = Cons 1 (Cons 2 (Var 2))
```

- Can have pointers outside of the list.
- This can be avoided by using dependent types:

Cyclic Lists as Nested Datatype

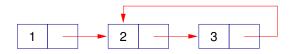
Definition:

- Var a represents a backward pointer to an element in a list.
- Nothing is the pointer to the first element of a cyclic list.
- Just Nothing is for the second element, and
- Just (Just Nothing) is for the third element, etc.
- The complete cyclic list has type CList Empty, where Empty is a type without constructors.

Examples



Cons 1 (Cons 2 (Var Nothing))



Cons 1 (Cons 2 (Cons 3 (Var (Just Nothing))))



Cons 1 (Cons 2 (Cons 3 Nil))

"Folding" a cycle

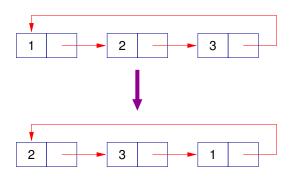
"Standard" fold

• Example:

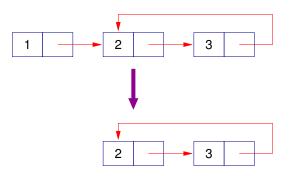
```
newtype K a = K Int summ = fold (\lambda x -> K 0) (K 0) (\lambda i (K j) -> K (i+j))
```

- In genereal, not very easy to use.
- Generalized folds for nested data types (Bird, Patterson)

Tail of a Cyclic List



Tail of a Cyclic List



Coalgebraic structure on Cyclic Lists

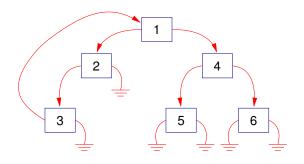
```
chead :: CList Empty -> Int
chead (Cons x 1) = x
ctail :: CList Empty -> CList Empty
ctail (Cons x 1) = csnoc x 1
csnoc :: Int -> CList (Maybe a) -> CList a
csnoc h Nil
                     = Nil
csnoc h (Var Nothing) = Cons h (Var Nothing)
csnoc h (Var (Just x)) = Var x
csnoc h (Cons x 1) = Cons x (csnoc h 1)
unwind :: CList Empty -> [Int]
unwind Nil = []
unwind 1 = chead 1 : unwind (ctail 1)
```

Cyclic Binary Trees

Definition:

- Nodes are "numbered" top-down.
- All nodes on the same level have the same "number".
- Has only backpointers to form cycles.
- Pointers to other directions forbidden, hence no sharing.

Example



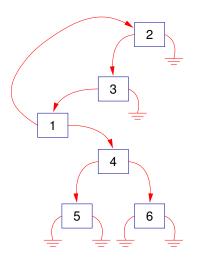
```
Node 1 (Node 2 (Node 3 (VarT Nothing) Leaf)

Leaf)

(Node 4 (Node 5 Leaf Leaf)

(Node 6 Leaf Leaf))
```

Cyclic children



Coalgebra structure on Cyclic Binary Trees

```
sonL :: CTree a -> CTree a
sonL (Node x t1 t2) = tsnocL x t2 t1
tsnocL :: Int -> CTree (Maybe a)
              -> CTree (Maybe a) -> CTree a
tsnocL x t (VarT Nothing) = Node x (VarT Nothing) t
tsnocL x t (VarT (Just n)) = VarT n
tsnocL x t Leaf
                          = Leaf
tsnocL x t (Node y t1 t2) = Node y (tsnocL x t' t1)
                                     (tsnocL x t', t2)
    where t' = relabel t
relabel :: CTree a -> CTree (Maybe a)
relabel (VarT x) = VarT (Just x)
relabel Leaf = Leaf
relabel (Node x t1 t2) = Node x (relabel t1) (relabel t2)
                                     4□ → 4□ → 4 □ → 1 □ → 9 Q (~)
```

Conclusions

- Generic framework to model cyclic structures.
- Backward pointers no sharing, just cycles.
- Type system guarantees the safety of pointers.
- To do: Work out all the details . . .
- New examples: zip, reverse, etc.