

SecreC for Sharemind 3

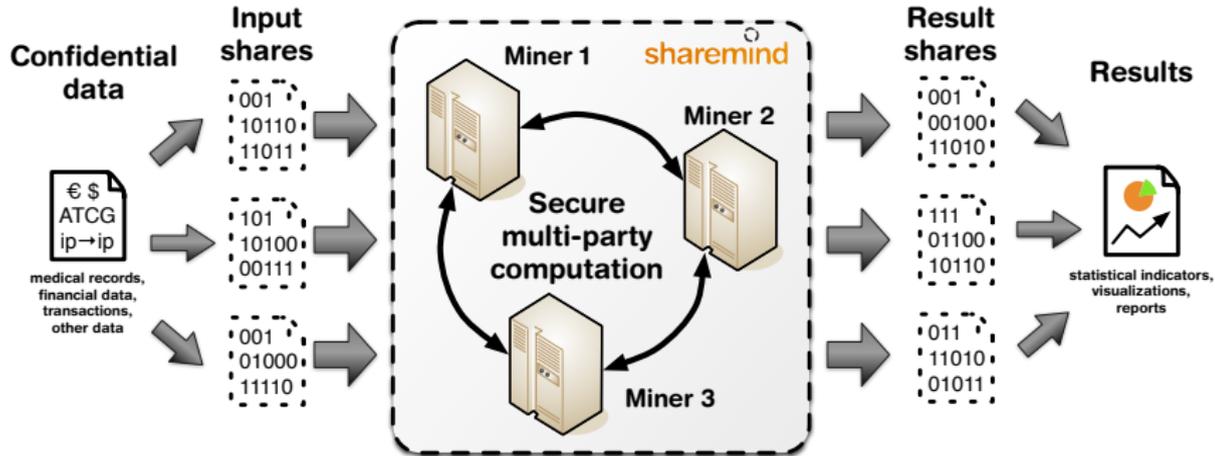
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Sharemind 2

- Sharemind is a secure multi-party computation framework.
- Sharemind 2 only supports additive 3-party secret sharing scheme.



SecreC for Sharemind 2

- Simple high-level imperative algorithm language.
- Hides implementation details of the secret sharing scheme.
- Two visibility types: public and private.
- Private values become public only through declassify.

```
void main () {                               // main function
    private uint a, b, c; // private data
    a = b + c;                               // private computation
    public uint d;                             // public data
    d = declassify (a); // private -> public
    publish (d);                               // send to client
}
```

Sharemind 3

- Various secure computation schemes.
- New underlying virtual machine.
- Complete rewrite of SecreC was in order.
- Some features of new SecreC:
 - more primitive data types
 - arbitrary dimensional arrays
 - simple module system
 - a lot of syntactic niceties
 - IR for data-flow analysis
 - protection domain polymorphism
 - etc.

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 - protection domain polymorphism
 - etc.

Protection domains

Definition (Protection domain)

A **protection domain** (PD) is a set of data that is protected with the same resources and for which there is a well-defined set of algorithms and protocols for computing on that data while keeping the protection.

Definition (Protection domain kind)

A **protection domain kind** (PDK) is a set of data representations, algorithms and protocols for storing and computing on protected data.

- Each PD belongs to a PDK, and each PDK can have several PD-s.

Protection domains

Example (protection domain kinds)

- A FHE system specified by its algorithms.
- A MPC system specified by its protocols.
- Public computation system.

Example (protection domains)

- A FHE system running under a single key.
- A single physical MPC instance.
- A public machine.

- Sharemind 2 supports a single protection domain called “private” in additive 3-party protection domain kind.

Protection domains in SecreC

- Simplest solution:
 - only public and private types
 - during program deployment map private to some protection domain
- Some issues:
 - impossible to use multiple PD-s concurrently
 - some code is PDK specific
- Better solution is to provide the ability to declare new protection domains.

Feature : Protection domains

```
kind additive3pp;           // declare PDK
domain pd_a3p additive3pp; // declare PD

void main () {              // main function
    pd_a3p uint a, b, c;     // private data
    a = b + c;               // private computation
    public uint d;           // public data
    d = declassify (a);      // private -> public
    publish (d);             // send to client
}
```

- declassify, and publish no longer primitives.
- Module declares the PDK.

Feature : Modules

```
import stdlib;           // import stdlib
import additive3pp;     // import PDK
domain pd_a3p additive3pp; // declare PD

void main () {          // main function
    pd_a3p uint a, b, c; // private data
    a = b + c;          // private computation
    public uint d;      // public data
    d = declassify (a); // private -> public
    publish (d);        // send to client
}
```

- Standard library declares publish.
- Module declares the PDK and declassify.

Feature : PD monomorphic functions (1/2)

```
kind additive3pp;  
domain pd_a3p additive3pp;  
  
pd_a3p uint sum (pd_a3p uint[[1]] x) {  
    pd_a3p uint s = 0;  
    for (uint i = 0; i < size (x); ++ i) {  
        s += x[i];  
    }  
  
    return s;  
}
```

Feature : PD monomorphic functions (2/2)

```
kind xor3pp;  
domain pd_x3p xor3pp;  
  
pd_x3p uint sum (pd_x3p uint[[1]] x) {  
    pd_x3p uint s = 0;  
    for (uint i = 0; i < size (x); ++ i) {  
        s += x[i];  
    }  
  
    return s;  
}
```

- Not very useful (apart from public PD).

Feature : PD polymorphic functions

```
template <domain D>
D uint sum (D uint[[1]] x) {
    D uint s = 0;
    for (uint i = 0; i < size (x); ++ i) {
        s += x[i];
    }

    return s;
}
```

- Type variable D binds to any PD or public.
- Implemented via code duplication.
- C++ templates.

Feature : overloading (1/2)

```
template <domain D : additive3pp>  
uint declassify (D uint x) {  
    // invoke system call that additive3pp  
    // is known to define  
}
```

```
template <domain D : xor3pp>  
uint declassify (D uint x) {  
    // invoke system call from xor3pp  
}
```

```
uint declassify (uint x) { return x; }
```

Feature : overloading (2/2)

```
template <domain D1 : additive3pp,  
          domain D2 : additive3pp>  
D1 uint reclassify (D2 uint x) {  
    // invoke a system call  
}
```

```
template <domain D>  
D uint reclassify (D uint x) { return x; }
```

- Overload selection via ad-hoc manner.
- For example: the number of instantiated PD variables.
- Operator overloading.

Feature : operator overloading

```
template <domain D : additive3pp>
D uint operator * (D uint x, D uint y) {
    // invoke system call
}
```

- Much of PDK functionality.

Expected use of polymorphism

Developer

- Standard library provides low-priority implementations.
- Each PDK is declared as a module.
 - declares the protection domain kind
 - PD polymorphic operations for that PDK
 - most operations through system calls

End user

- Algorithm is PD polymorphic.
- Main program fixes the PD.

Example : sorting (1/2)

```
module stdlib;

template <domain D>
D uint[[1]] sort (D uint[[1]] src) {
    uint[[2]] sn = sortnet (size (src));
    for (uint i = 0; i < shape (sn)[0]; ++ i) {
        D uint x = src[sn[i,0]];
        D uint y = src[sn[i,1]];
        D uint b = (uint) (x < y);
        src[sn[i,0]] = x*b + y*(1 - b);
        src[sn[i,1]] = x*(1 - b) + y*b;
    }

    return src;
}
```

Example : sorting (2/2)

```
module additive3pp;

template <domain D : additive3pp>
D uint [[1]] sort (D uint [[1]] src) {

    src = shuffle (src);

    // Sort using:
    //   declassify (src[i] < src[j])

    return src;
}
```

- Can be done in $O(n \log n)$.

Example : multiple protection domains

```
template <domain D1, domain D2>
D1 uint hamming (D2 uint[[1]] x, D2 uint[[1]] y) {
    D2 bool[[1]] neqs = (x != y);
    D1 uint[[1]] vs = (uint) reclassify (neqs);
    return sum (vs);
}
```

- Different performance depending on selected PD-s.
- If D1 = pd_a3p, and D2 = pd_a3p then 12 rounds.
- If D1 = pd_a3p, and D2 = pd_x3p then 8-9 rounds.

Summary

Good

- Simple solution.
- Easy to understand.

Bad

- Ad-hoc overload resolution is not intuitive.
- Bad type errors just like in C++.
- Template interaction with modules can be strange.
 - specialization could help

Semantics

- Formally defined type system.
- Small-step operational trace semantics.
 - labels are system calls
- Monomorphic intermediate language.
- Type-directed translation from polymorphic to monomorphic language.
- Weak bi-simulation between semantics of monomorphic and polymorphic language.
- Security of information flow.
- Models the compiler with a monomorphic IR.
- Abstract syntax, type system and semantics have been invaluable documentation for implementing and debugging the compiler.

Future work

- Language improvements.
 - pass-by-reference
 - user defined data types
 - many small improvements
- Standard library.
- We need users and more PD-s.
- Something better than templates?