

## Konstantidega lambda-arvutus

- Konstantidega  $\lambda$ -termide süntaks:

$$\begin{array}{l}
 e ::= x \\
 \quad | c \\
 \quad | (e_1 e_2) \\
 \quad | (\lambda x. e)
 \end{array}
 \qquad
 \begin{array}{l}
 \text{muutuja} \\
 \text{konstant} \\
 \text{aplikatsioon} \\
 \text{abstraktsioon}
 \end{array}$$

- Iga konstandiga seotakse mingi arv  $\delta$ -reduktsiooni reegleid.
- Näide: Naturaalarvud  $(0, 1, 2, \dots)$  ja liitmine  $(+)$ .

$$\begin{array}{l}
 + 0 0 \rightarrow_{\delta} 0 \\
 + 0 1 \rightarrow_{\delta} 1 \\
 \dots
 \end{array}
 \qquad
 \begin{array}{l}
 + 1 0 \rightarrow_{\delta} 1 \\
 + 1 1 \rightarrow_{\delta} 2 \\
 \dots
 \end{array}$$

- $\delta$ -reeglite lisamine võib kokkuvoolavuse ära rikkuda!

## Andmete esitamise $\lambda$ -arvutuses

- Baaskombinaatorid

$$I \equiv \lambda x. x$$

$$K \equiv \lambda x y. x$$

$$S \equiv \lambda f g x. f x (g x)$$

- “Astendamine”

$$\begin{aligned}
 E^0 E' &\equiv E' \\
 E^n E' &\equiv \underbrace{E(E(\dots(E E')\dots))}_{n \text{ tükki}}
 \end{aligned}$$

- NB!**

$$E^n(E E') \equiv E^{n+1} E' \equiv E (E^n E')$$

## Tõeväärtused

- Spetsifikatsioon

not true = false

not false = true

- Definiitsioon

true  $\equiv \lambda xy. x$  ( $\equiv K$ )

false  $\equiv \lambda xy. y$

not  $\equiv \lambda t. t \text{ false true}$

- Näide:

not true  $\equiv (\lambda t. t \text{ false true}) \text{ true}$

$\rightarrow \text{true false true}$

$\equiv (\lambda x. \lambda y. x) \text{ false true}$

$\rightarrow (\lambda y. \text{false}) \text{ true}$

$\rightarrow \text{false}$

## Tingimuslause

- Spetsifikatsioon

$$\begin{aligned} \text{cond true } E_1 E_2 &= E_1 \\ \text{cond false } E_1 E_2 &= E_2 \end{aligned}$$

- Definiatsioon

$$\text{cond} \equiv \lambda t x y. t x y$$

- Näide:

$$\begin{aligned} \text{cond false } E_1 E_2 &\equiv (\lambda t x y. t x y) \text{ false } E_1 E_2 \\ &\rightarrow \text{false } E_1 E_2 \\ &\equiv (\lambda x y. y) E_1 E_2 \\ &\rightarrow E_2 \end{aligned}$$

## Paarid ja ennikud

## • Paarid

$$\begin{aligned} \text{fst} &\equiv \lambda p. p \text{ true} \\ \text{snd} &\equiv \lambda p. p \text{ false} \\ (E_1, E_2) &\equiv \lambda f. f E_1 E_2 \end{aligned}$$

## • Ennikud

$$\begin{aligned} (E_1, \dots, E_n) &\equiv (E_1, (\dots (E_{n-1}, E_n) \dots)) \\ E \downarrow^n 1 &\equiv \text{fst } E \\ E \downarrow^n 2 &\equiv \text{fst } (\text{snd } E) \\ &\dots \\ E \downarrow^n i &\equiv \text{fst } (\text{snd}^{i-1} E) \\ &\dots \\ E \downarrow^n n &\equiv \text{snd}^{n-1} E \end{aligned}$$

## Naturaalarvud

- Standardnumbrid

$$\begin{aligned} \ulcorner 0 \urcorner &\equiv \lambda x. x && (\equiv I) \\ \ulcorner n+1 \urcorner &\equiv (\text{false}, \ulcorner n \urcorner) \\ \text{succ} &\equiv \lambda n. (\text{false}, n) \\ \text{pred} &\equiv \lambda n. n \text{ false} && (\equiv \text{snd}) \\ \text{iszero} &\equiv \lambda n. n \text{ true} && (\equiv \text{fst}) \end{aligned}$$

- Liitmine (!?)

$$\text{add} = \lambda x y. \text{cond}(\text{iszero } x) y (\text{add}(\text{pred } x)(\text{succ } y))$$