Quantum Cryptography (spring 2021)

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## Exercise Sheet 13

Out: 2021-05-18

Due: 2021-05-25

This is a **bonus homework**. It can be used if you have less than 50% of the point total. (Recall that you need 50% to qualify for the exam.) But you are also welcome (and encouraged) to solve the homework if you do not need the points, for sake of learning.

## 1 Quantum proofs

Knowlets:	ProofSys	ProblemID: QProofs
Time:		
Difficulty:		

Show that if (P, V) is a proof system (Definition 53 in the lecture notes), then it also is a quantum proof system as in the following definition:

**Definition 1 (Quantum proof systems)** We call a pair (P, V) of interactive machines a quantum proof system for the relation R with soundness-error  $\varepsilon$  iff the following two conditions are fulfilled:

- Completeness: For any  $(x, w) \in R$ , we have that  $\Pr[\langle P(x, w), V(x) \rangle = 1] = 1$ .
- Soundness: For any (potentially computationally unlimited) quantum machine  $P^*$ , and for any  $x \notin L_R$ , we have  $\Pr[\langle P^*(), V(x) \rangle = 1] \leq \varepsilon$ .

Notice that the only difference to Definition 53 in the lecture notes is the additional word **quantum**.

## 2 Zero-knowledge and discrete logarithm

Fix a group G of prime order q with generator g.  $(G, q, \text{ and } g \text{ may depend on some implicit security parameter but are considered publicly known.) Let <math>R := \{(x, w) : g^w = x, w \in \{0, \dots, q-1\}\}.$ 

Consider the following proof system for R (Schnorr's proof system for discrete logarithms):

- The prover P gets input  $(x, w) \in R$ .
- The verifier V gets input  $x \in R$ .
- The prover P chooses  $b \stackrel{\$}{\leftarrow} \{0, \ldots, q-1\}$  and sends  $a := g^b$  to the verifier V.
- The verifier chooses  $r \stackrel{\$}{\leftarrow} \{0, \ldots, q-1\}$  and sends r to the prover P.
- The prover P computes  $s := b + rw \mod q$  and sends s to the verifier V.
- The verifier V checks whether  $x, a \in G$  and  $g^s = ax^r$ .

This proof system is well-known to be a proof system. However, in the classical setting, it is unknown whether this proof system is zero-knowledge!<sup>1</sup>

	Knowlets:	ProofSys	ProblemID: ZKDlogSound
(a)	Time:		
	Difficulty:		
	Show that $(P, V)$ is a proof system with soundness-error $1/q$ .		

	Knowlets:	QZK, DlogAlgo	ProblemID: ZKDlogShor
(b)	Time:		
	Difficulty:		

Show that (P, V) is statistical quantum zero-knowledge.

Hint: This has nothing to do with rewinding! It has a lot to do with Shor's algorithm.

<sup>&</sup>lt;sup>1</sup>It is however "honest-verifier zero-knowledge". This is a weaker notion where the verifier is considered to behave honestly.