

Program Transformation with Stratego/XT

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Based on
Program Transformation with Stratego/XT.
Rules, Strategies, Tools, and Systems in
Stratego/XT 0.9.
by Eelco Visser, 2004

What Is Program Transformation?

- Automatic manipulation of source programs
- Used in
 - Compilers
 - Code generators
 - Refactoring tools
 - Migrating tools
 - Reverse engineering tools

What Is Stratego/XT?

- Stratego is a language for
 - Transformation rules
 - Programmable strategies for applying these rules
- XT is a toolset containing
 - Parser generators
 - Pretty-printer generators
 - Grammar engineering tools

Transformation Rule

- Encodes a basic transformation step as a rewrite on an abstract syntax tree
- Examples:
 - EvalPlus : $\text{Plus}(\text{Int}(i), \text{Int}(j)) \rightarrow \text{Int}(k)$
where $\langle \text{add} \rangle(i, j) \Rightarrow k$
 - LetSplit : $\text{Let}([d1, d2 \mid d^*], e^*) \rightarrow \text{Let}([d1], \text{Let}([d2 \mid d^*], e^*))$

Representing Terms

- Rewrite rules apply to abstract syntax trees
 - As opposed to parse trees
- Abstract syntax trees are represented by first-order prefix terms
 - E.g., `If(cond, then, else)`
- Syntax tree fragments can be described using the concrete syntax of the object language.
 - `EvalPlus : |[i + j]| -> |[k]|`
 where `<add>(i, j) => k`

Transformation Rule (2)

- Rewrite rules can be broken down into the more basic actions
 - Matching
 - Building
 - Variable scope
- Typically, rewrite rules are context-free
 - Scoped dynamic rewrite rules can be used to generate rewrite rules in runtime
 - Dynamic rules can encode contextual information

Term Rewriting

- Term rewriting is the exhaustive application of a set of rewrite rules to a term until no rule applies anywhere in the term
 - Also called normalization
- Example
 - `Minus (Plus (Int (4), Plus (Int (1), Int (2))), Var ("a"))`
 - ... reduces to `Minus (Int (7), Var ("a"))` by repeatedly applying **EvalPlus**

Term Rewriting (2)

- Exhaustive rewriting is used in most rewriting tools
- However, normalizing a term with respect to *all* rules is not always desirable
- Often this is done by using special kind of rules
 - Example: encode evaluation order in rules

Programmable Rewriting

- Stratego makes rewriting strategy *explicit* and *programmable*
- One has to define explicitly
 - Which rules to apply
 - Which strategy to follow
- Example
 - `simplify = innermost(EvalPlus + LetSplit + ...)`

Transformation Strategy

- *Strategy* is an algorithm that transforms a term into another term or fails at doing so
- Combines a set of rules into a complete transformation
 - Orders their application using control and traversal combinators
- Important property: ability to define generic traversals
 - Do not depend on specific data types

Strategy Combinators

- Stratego's approach is to allow building complex strategies from very simple building blocks
 - Sequential composition ($s1 ; s2$)
 - Deterministic choice ($s1 <+ s2$; first try $s1$, only if that fails $s2$)
 - Non-deterministic choice ($s1 + s2$; same as $<+$, but the order of trying is not defined)
 - Guarded choice ($s1 < s2 + s3$; if $s1$ succeeds then commit to $s2$ else $s3$)

Strategy Combinators (2)

- Building blocks
 - Testing (where(s); ignores the transformation achieved)
 - Negation (not(s); succeeds if s fails)
 - Recursion (rec x(s))

Strategy Definitions

- $f(x_1, \dots, x_n) = s$
 - Define strategy f
 - $x_1 \dots x_n$: strategy arguments
- Examples
 - `try(s) = s <+ id`
Applies strategy s , succeeds even if it fails.
 - `repeat(s) = try(s; repeat(s))`
Repeats transformation s until it fails

Strategy Definitions (2)

- Strategy definitions do not explicitly mention the term to which they are applied
- Instead, they combine term transformations into more complex term transformations

Congruence Operator

- Basically match and apply

- `control-flow(s) =`
 `Assign(id, s)`
 `+ If(s, id, id)`
 `+ While(s, id)`
 - `map(s) = [] + [s | map(s)]`

- Defines traversals that are specific to a data type

Generic Traversals

- Not specific to any data type
- One-pass traversals:
 - `all(s)` – applies `s` to each subterm of current term
 - `bottomup(s) = all(bottomup(s)); s`
 - `topdown(s) = s; all(topdown(s))`
 - `alltd(s) = s <+ all(alltd(s))`
 - `oncetd(s) = s <+ one(oncetd(s))`

Generic Traversals (2)

- **Fixpoint traversals:**

- `innermost(s) =
 bottomup(
 try(s; innermost(s))`

Cascading Transformations

- Applying several small transformation steps.
 - `simplify =`
`innermost (R1 <+ ... <+ Rn)`

Staged Transformations

- Transformations are applied in stages
 - `simplify =`
 `innermost(A1 <+ ... <+ Ak)`
 `; innermost(B1 <+ ... <+ B1)`
 `; ...`
 `; innermost(C1 <+ ... <+ Cm)`
- Can be combined with cascading transformations
- Because rules are independent from strategies, they can be reused in different stages

Local Transformations

- Transformations are applied only to subtree of the program where they make the most sense
 - `transformation =
 alltd(
 trigger-transformation
 ; innermost (A1 <+ ... <+ An)
)`
 - *trigger-transformation* selects one node where the cascading transformation is applied

First-Class Pattern Matching

- In addition to rules, pattern matching is available as primitives in strategies
 - `EvalPlus : Plus(Int(i), Int(j)) -> Int(k)`
 `where <add> (i, j) => k`
 - `EvalPlus = {i,j,k: ?Plus(Int(i), Int(j));`
 `where(! (i,j); add; ?k); !Int(k) }`

Dynamic Rules

- Pure rewriting rules and strategies are context-free
- Passing context through arguments can quickly become tedious
 - List of defined functions
 - List of defined variables
- Stratego offers mutable global state in the form of dynamic rules

Dynamic Rules (2)

- Generated at run-time
- Can access information available from their generation context

Dynamic Rule Example

```
DeclareFun =  
  ?fdec@[ [ function f(x1*) ta = e1 ] ] ;  
  rules(  
    InlineFun :  
      |[ f(a*) ]| -> |[ let d* in e2 end ]|  
      where <rename> fdec =>  
          |[ function f(x2*) ta = e2 ]|  
          ; <zip(BindVar)>(x2*, a*) => d*  
    )  
  BindVar :  
    (FArg |[ x ta ]|, e) ->  
    |[ var x ta := e ]|
```


Dynamic Rule Scope

- Restricts the scope of new definitions of the dynamic rule
 - Rule is removed if execution goes out of scope
- Example:
 - ```
inline = {| InlineFun:
 try(DeclareFun)
 ; repeat(InlineFun + Simplify)
 ; all(inline)
 ; repeat(Simplify)
 |}
```

# Term Annotations

- Abstract syntax of programs is expressed in terms
  - Term = Constructor + list of argument terms
- Sometimes it is useful to attach information to term without changing it
- This information can be stored in annotations
  - Each term has list of annotations
  - Annotations are also terms

# Term Annotations (2)

- Annotations can be processed like any other terms
- Example:
  - `TypeCheck : Plus(e1{Int}, e2{Int}) -> Plus(e1, e2){Int}`
- Can be used for
  - Type checking
  - Strictness analysis
  - Bound-unbound variables analysis

# Problems with Annotations

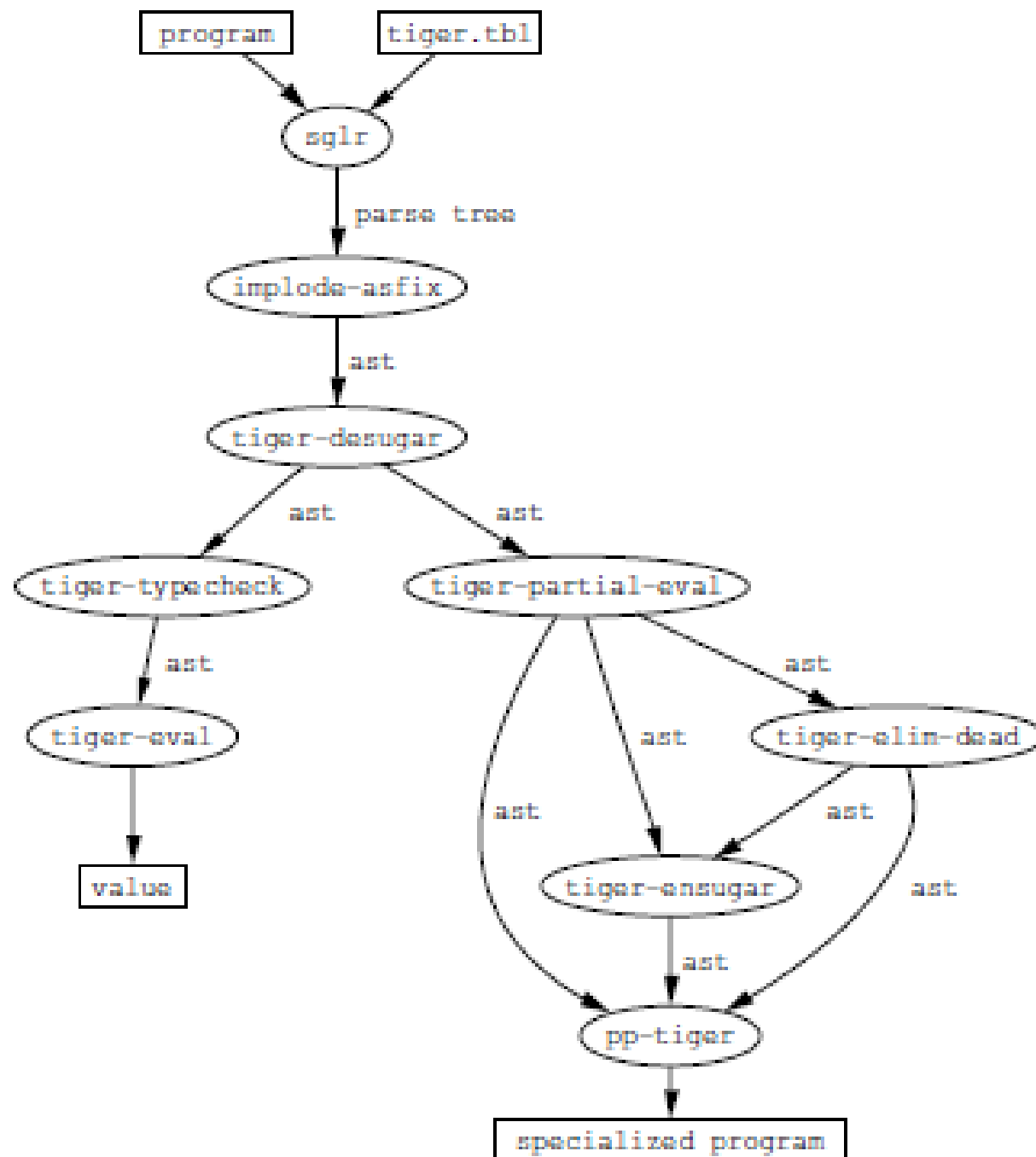
- Transformations are supposed to preserve annotations
  - Stratego's traversals do that
  - When transforming a term, this is not so simple, because this should follow semantics of the annotations
- Should annotations affect equality between terms?

# Transformation Tool

- Wraps a composition of rules and strategies into a stand-alone, deployable component
- Can be called from the command-line or from other tools
- Transforms terms into terms
  - All the tools in Stratego/XT toolkit use standard ATerm format for terms

# Transformation System

- Composition of tools that performs a complete source-to-source transformation
- Consists of:
  - Parser
  - Pretty-printer
  - Transformation tools





Thank You